



2021 Australasia **Db2** Tech Conference

Db2 reorg unleash the full power of the utility

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#### Agenda

- Basics (why and how to reorg)
- Classic vs inplace
- How to run a reorg in parallel
- Alternative to reorg (amt)
- Performance Comparison



#### Reorgs are recommended, because (1|2)

- Someone executed a certain alter table command
- The following is the full list of REORG-recommended ALTER statements that cause a version change and place the table into a REORG-pending state:
  - DROP COLUMN
  - ALTER COLUMN SET NOT NULL
  - ALTER COLUMN DROP NOT NULL
  - ALTER COLUMN SET DATA TYPE, except in the following situations:
    - Increasing the length of a VARCHAR or VARGRAPHIC column
    - Decreasing the length of a VARCHAR or VARGRAPHIC column without truncating trailing blanks from existing data, when no indexes exist on the column
- Two (or more) alter commands in one UOW increase the counter by 1



#### Reorgs are recommended, because (2 | 2)

- up to 3 single reorg recommended operations allowed
- some commands possible if table in reorg pending state:
  - Drop table
  - Rename table
  - Truncate table
  - Reorg table (offline and full table only)
- Check for current counters: SYSIBMADM.ADMINTABINFO (use a few columns and a where clause to filter for specific table(s) or schema)



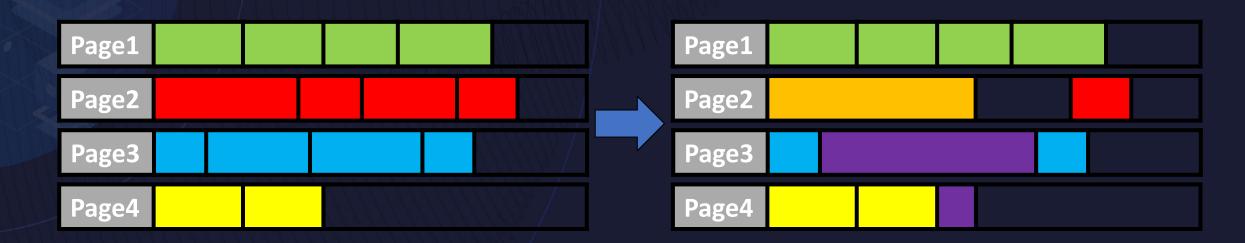
#### Reorgs are needed (reorg pending), because

- Reorg recommended counter is >0 / reorg pending state
- Regulare tablespace has reached max pages limit and has to be converted to large (you will need a reorg indexes to "activate" large RIDs and make table "insertable" again)



#### Reorgs are really recommended, because (1|2)

- Reorgcheck shows at least 2\* on a table or 3\* on an index (data is really bad fragmented)
- Performance is slow (having overflow records), like after adding new columns or heavy updates on variable columns with larger data





#### Reorgs are really recommended, because (2 | 2)

- compression is now inplace (generate compression dictionary)
  - Generate a dictionary because of high amount of new data that is compressed badly
  - After a load operation as records used to generate initial dictionary are a bad sample
  - After changing from dictionary compression to adaptiv compression
- Everytime you convert tablespace from regular to large



#### how execute a reorg (1|3)

- via command line
  - db2 "reorg table myschema.mytable"
- Via SQL
  - CALL SYSPROC.ADMIN\_CMD ('REORG TABLE myschema.mytable')
  - CALL SYSPROC.ADMIN\_REVALIDATE\_DB\_OBJECTS('table', 'MYSCHEMA', 'MYTABLE')
     only for tables in reorg pending state, will execute and runstats for those



#### how execute a reorg (2 | 3)

- Can be executed for a table
  - db2 "reorg table myschema.mytable"
- Can be executed for a data partition of a table (range parted)
  - db2 "reorg table myschema.mytable on data partition part001"
- Can be executed for a database partition of a table (DPF)
  - db2 "reorg table myschema.mytable on DBPARTITIONNUM (1)"



#### how execute a reorg (3 | 3)

- Can be executed for all indexes of a table
  - db2 "reorg indexes all for table myschema.mytable"
- Can be executed for a given index of a table
  - db2 "reorg index myindex for table myschema.mytable"
  - Only supported for:
    - Nonpartitioned indexes on a data partitioned table that are not block indexes
    - Any index on any permanent table if CLEANUP ALL is specified and RECLAIM EXTENTS is not specified



#### Pitfalls when executing a reorg (1|3)

- Table reorg always recreates all indexes as well
  - Maybe better to drop idx / reorg table / create idx (especially when reorg gets logged in hadr -> avoid log full)
- Reorg a table partion with non parted indexes will lead to a reorg of the non parted idx for the whole table
- If partitioned table is in reorg pending and is having non partitioned indexes you will have to reorg whole table at once (not per partition)
- Inplace / online reorg is not allowed when in reorg pending state



#### Pitfalls when executing a reorg (2|3)

- By default lobs are not reorganized (use parameter LONGLOBDATA)
- By default db2 uses the same tablespace for the working copy (if partitioned it will use the tablespace of that partition)
  - At least twice space is needed in the same tablespace or specify a system temp space to write working copy in that tablespace
- Reorg on the complete partitioned table will result in 1 partition reorg at a time (at least twice space if largest partition is needed)
- "The original table might be available for queries until the replace operation starts, depending on the access clause" (classic/offline reorg)



#### Pitfalls when executing a reorg (3|3)

- Allow read access on partitioned table only supported when having no non partitioned indexes (otherwise it will be allow no access) - all other partitions would be read and writeable
- When having a non partitioned index the whole table will be set in no access mode (not only the given partition)
- By default no automatic parallelism for a table (it will always use a single thread, one partition after another)



#### How to avoid reorgs

- Use more alter table statements in one OUW to avoid getting from reorg recommended to reorg pending status
- Minimize the usage of VARCHAR, VARGRAPHIG to avoid overflows
- Use ITC (Insert Time Clustering) tables data and indexes will be stored and deleted in blocks - no reorg needed (and not supported) and having a new not usable column and index to store the insert time
- Use admin move table before it comes to reorg pending situation



#### Monitoring a reorg

- GET SNAPSHOT FOR TABLES command
  - db2 get snapshot for tables on sample
- db2pd -reorg command
  - db2pd -db sample -reorg -repeat 60
- List History Reorg:
  - db2 list history reorg all for sample
- Select from the SYSIBMADM.SNAPTAB\_REORG
- INPLACE\_REORG\_STATUS in SYSPROC.ADMIN\_GET\_TAB\_INFO
- Use third party monitoring tool



#### Classic vs. inplace (offline vs. online) (1 | 2)

- Classic
  - "No" access to the table
  - Uses a copy approach, building a full copy of the table
  - Fast
  - Perfectly clustered data once finished
  - Indexes gets rebuild automatically
  - Can use temp tablespace to reduce needed storage in tablespace

- Inplace
  - Full access to the table
  - Rows are moved within the table, sequentielly
  - Slow
  - Maybe inperfect clustering depending on sql during reorg
  - Indexes maintained but not rebuild
  - Low storage requirement as using source object to move data



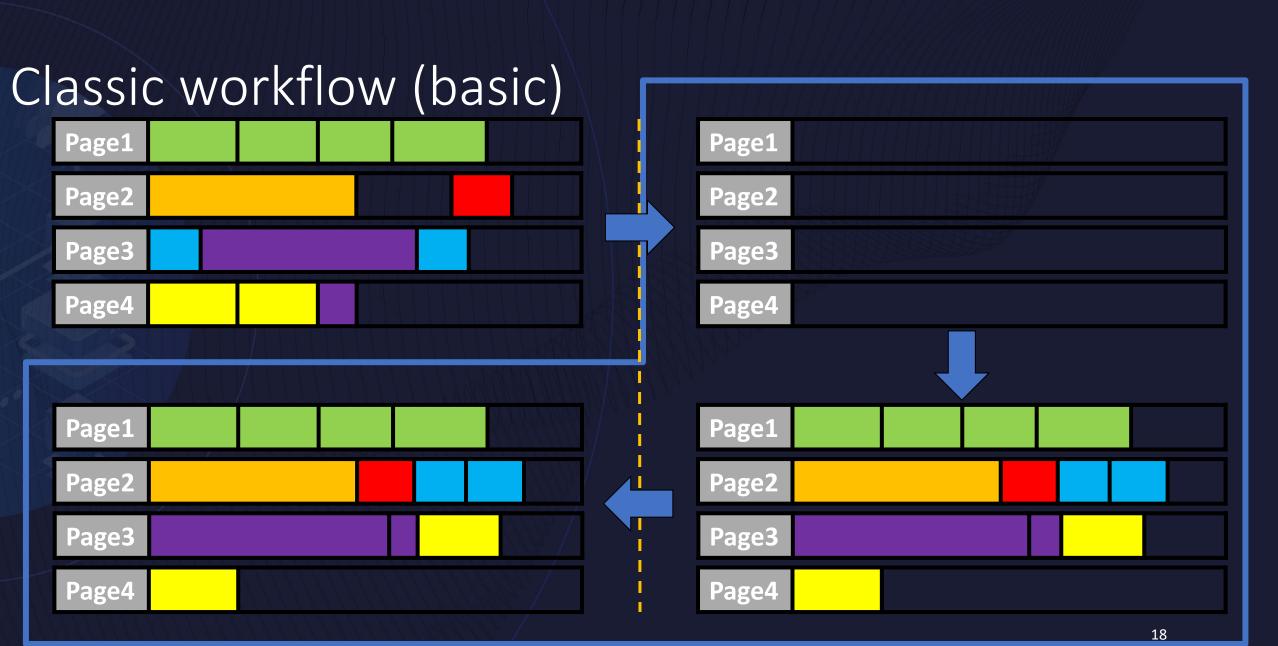
#### Classic vs. inplace (offline vs. online) (2 | 2)

- Classic
  - Large space requirement (shadow copy)
  - No control during execution (has to be restarted if it stops)
  - Not recoverable
  - Needs fewer TRX Log
  - Benefit once the reorg completes
  - Suitable for large tables

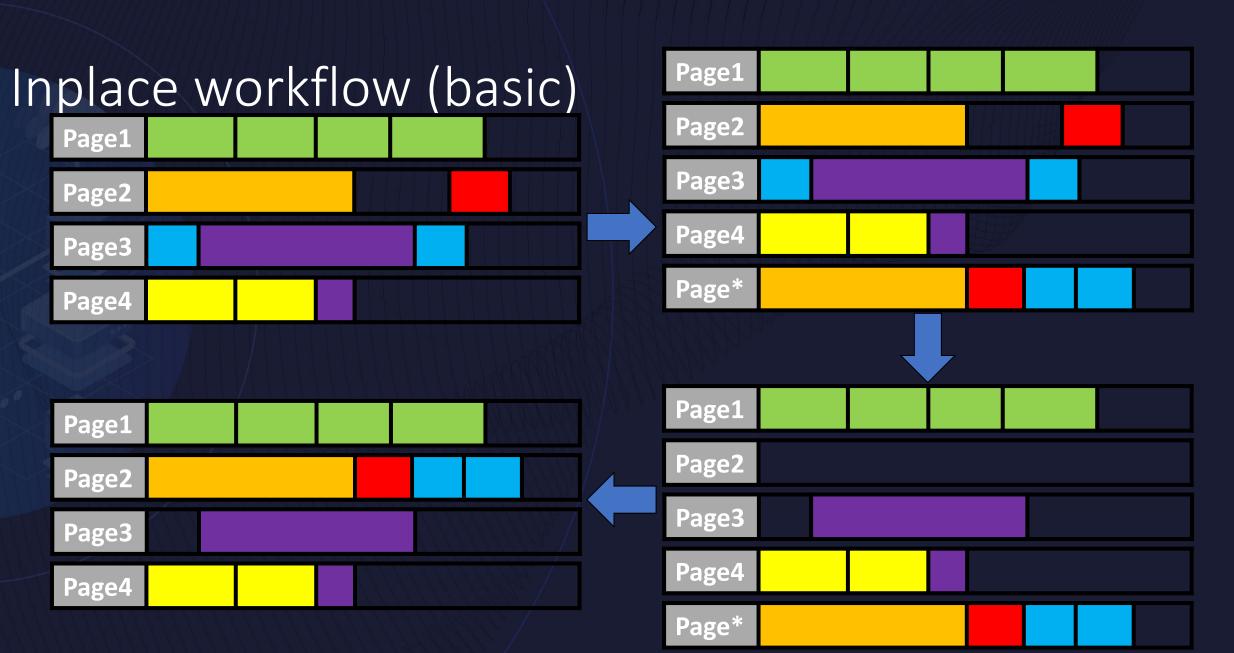
#### Inplace

- Small space requirement (moving data in existing object)
- Good control during the execution (pause, resume, start, stop)
- Fully recoverable
- Needs more TRX Log
- Direct benefit right after reorg starts
- Not suitable for large tables (may never finish)











#### Reorg options for partitioned Tables (1 | 4)

- A reorg can run on a single partition of range partitioned table, allowing access on all other partitions during the execution
- Multiple partition reorgs can run in parallel with following criteria:
  - Each REORG command must specify a different partition with the ON DATA PARTITION clause.
  - Each REORG command must use the ALLOW NO ACCESS mode to restrict access to the data partitions.
  - The partitioned table must have only partitioned indexes if issuing REORG
     TABLE commands. No nonpartitioned indexes (except system-generated XML path indexes) can be defined on the table.



#### Reorg options for partitioned Tables (2 | 4)

- Restrictions for an index to be partitionable
  - To be partitioned the index will have to include the distribution key create index parted\_1 on myschema.mytable (distri\_col,col1) partitioned
  - If distribution key is/can not be included, the index can be partitioned (if index is not unique) and the column will be included (hidden) -> maybe this index will be used differently in SQL queries (Data-partitioned secondary index (DPSI) on z/OS) create index parted\_2 on myschema.mytable (col1) partitioned



#### Reorg options for partitioned Tables (3 | 4)

- Examples:
  - Table is partitioned by quarters of a year, reorg per quarter or in parallel
    - REORG TABLE myschema.mytable ALLOW NO ACCESS ON DATA PARTITION P1
    - REORG TABLE myschema.mytable ALLOW NO ACCESS ON DATA PARTITION P2
    - REORG TABLE myschema.mytable ALLOW NO ACCESS ON DATA PARTITION P3
    - REORG TABLE myschema.mytable ALLOW NO ACCESS ON DATA PARTITION P4
  - What also is possible to be excuted in parallel
    - REORG INDEXES ALL FOR TABLE myschema.mytable ALLOW NO ACCESS ON DATA PARTITION P1
    - REORG TABLE myschema.mytable ALLOW NO ACCESS ON DATA PARTITION P2
    - REORG INDEXES ALL FOR TABLE myschema.mytable ALLOW NO ACCESS ON DATA PARTITION P3



#### Reorg options for partitioned Tables (4 | 4)

- Recommendations:
  - Use separate/dedicated tablespaces for every partition or group partitions if having several disks/volumes per path
  - Parallelism should not be higher than number of (free) Cores and number of disks/Volumes (one reorg will use full power of one core and underlying storage)





#### Admin Move Table - Basics (1|3)

- System Procedure officially introduced in v9.7
- Basically developed to move on table from one tablespace to another without blocking the table for UID
- Easy to use but flexibel for advanced usage (one command for the complete workflow vs. one command per step and more customization options)
- Columns can be added/dropped -> minize risk of reorg pending
- Improvements over last versions, getting less restrictions and more options



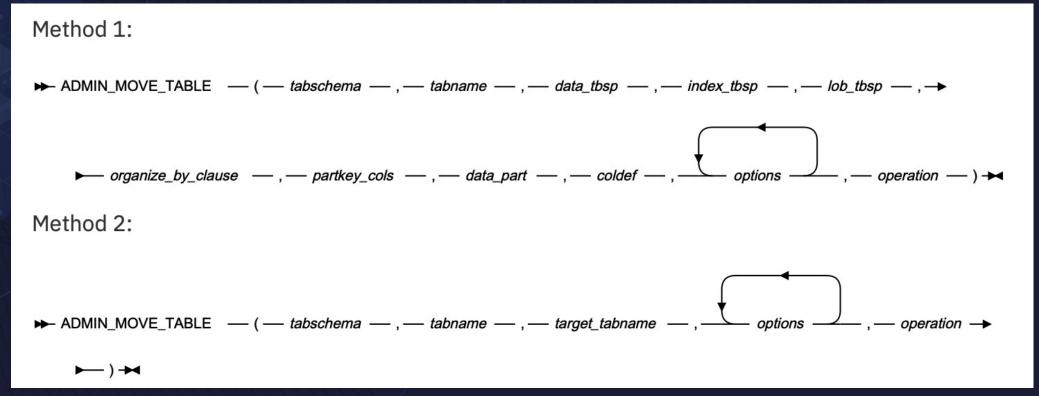
#### Admin Move Table - Basics (2 | 3)

- Admin Move Table has basically 4 needed steps:
  - INIT: Verifies requirements, creates target and staging tables and triggers
  - COPY: Copies data from source to target table (Insert or Load)
  - REPLAY: Move Data from staging to target (repeat this step if first execution took long time)
  - SWAP: Rename the target table to source table, rename indexes, remove staging table, remove source table (optionally keep it)
- And just one to go complete through, to terminate and to clean:
  - Move: Performs INIT, COPY, REPLAY, and SWAP operations in one step
  - Term: Terminates a running or killed table move
  - CANCEL: clears up all intermediate data



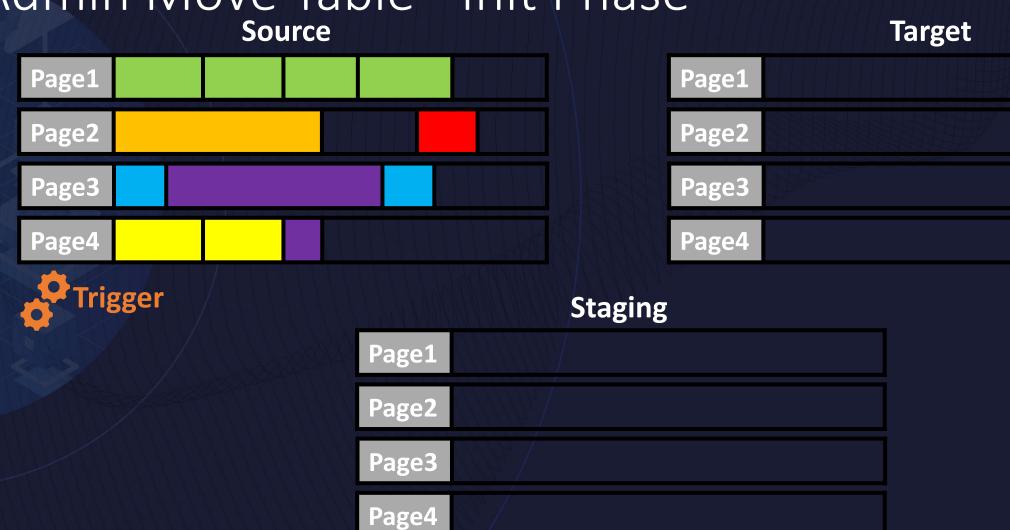
#### Admin Move Table - Basics (3|3)

 Two Procedure Methods (easy and more options, non-optional: tabschema, tabname and operation)



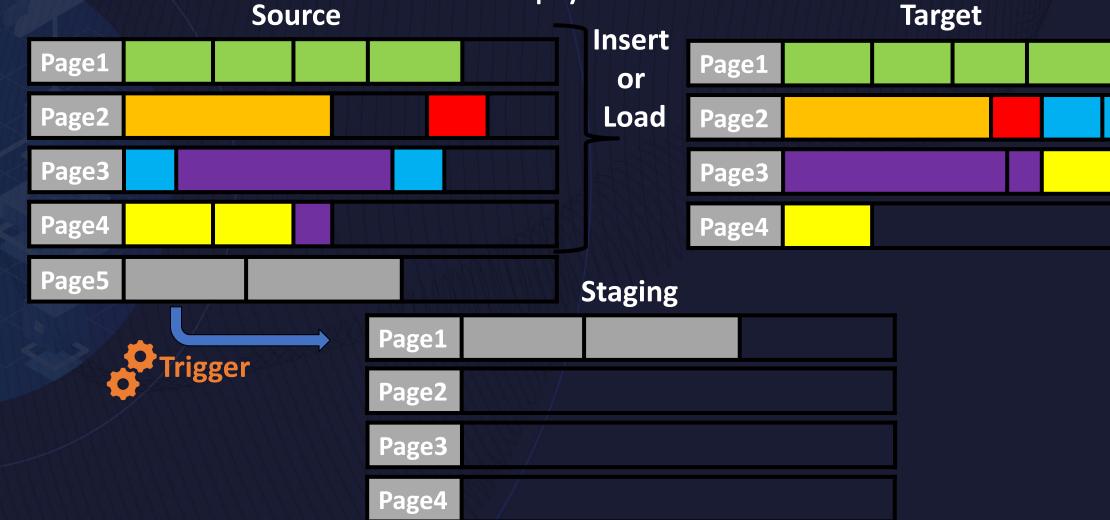


#### Admin Move Table - Init Phase

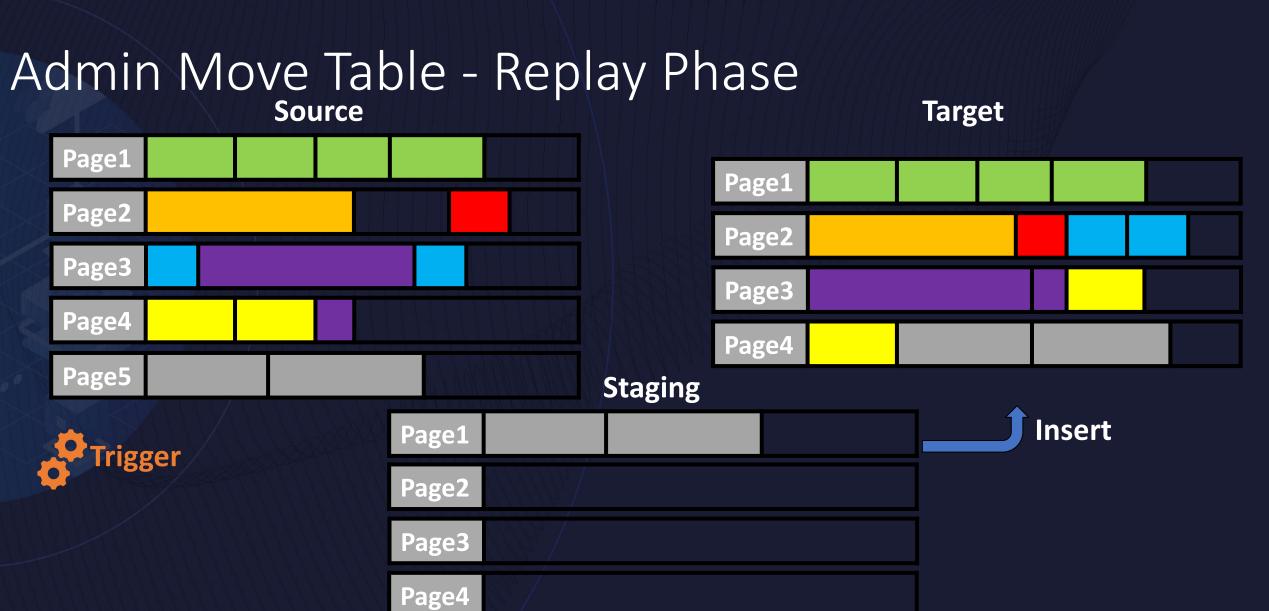




# Admin Move Table - Copy Phase Source









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#### Admin Move Table - alternative to reorg

- Can only be used for "really" reorg recommended not for reorg pending tables (use it before having reorg pending status)
- More or less fully online during swap phase X lock needed on source, rebinds needed after the swap
- Needs at least twice the space as source (maybe less if compression rate improves or had a lot of gaps in source)
- Reads data from source by default with order by if a clustering index exist or cluster option is specified (and clustering index, uniq index or primary key exists or an alternative index to use has been specified with ADMIN\_MOVE\_TABLE\_UTIL)



#### Admin Move Table - with an extra reorg

- Admin Move Table has the option to explicitly perform a reorg on target table during the swap step right before swap
- Only method to get an optimal compression dictionary for xml columns
- Set the REORG option at any point up to and including the SWAP phase
- Not recommended for "normal" or every day use (e.g. to compact/compress data/remove gaps/sort etc. it is not needed)



#### Admin Move Table - example

```
[[db2inst1@db2_test ~]$ db2 "call SYSPROC.ADMIN_MOVE_TABLE('TESTSCHEMA','TESTTABLE','','','MOVE')";
  Result set 1
  KEY
                                   VALUE
  AUTHID
                                   DB2INST1
  CLEANUP_END
                                   2020-07-30-12.30.38.146171
  CLEANUP_START
                                   2020-07-30-12.30.38.055724
 COPY_INDEXNAME
                                   CLUS_IDX
                                   TESTSCHEMA
 COPY_INDEXSCHEMA
  COPY_OPTS
                                   ARRAY_INSERT,CLUSTER_OVER_INDEX
  COPY_TOTAL_ROWS
  DICTIONARY_CREATION_TOTAL_TIME
                                   CLUS_IDX
  INDEXNAME
  INDEXSCHEMA
                                   TESTSCHEMA
  INDEX_CREATION_TOTAL_TIME
  INIT_END
                                   2020-07-30-12.30.37.185235
                                   2020-07-30-12.30.36.535411
  INIT_START
  ORIGINAL_TBLSIZE
                                   512
  REPLAY_END
                                   2020-07-30-12.30.37.986532
  REPLAY_START
                                   2020-07-30-12.30.37.518894
  REPLAY_TOTAL_ROWS
  REPLAY_TOTAL_TIME
  STATUS
                                   COMPLETE
  SWAP_END
                                   2020-07-30-12.30.38.040619
  SWAP_RETRIES
  SWAP_START
                                   2020-07-30-12.30.38.018785
  UTILITY_INVOCATION_ID
                                   0100000010000008000000000000000002020073012303718769600000000
  VERSION
                                   11.05.0400
```



#### Admin Move Table - to avoid reorg pending (1|3)

- Use admin move table for your DDL changes (change the target)
  - Drop / add column
  - Change data type (when types are compatible and column name persists)
  - Change nullable options
  - Primary Key
- Option1
  - Create your target table in forehand with structure of next version (only table!)
- Option2
  - Use parameter coldef of admin move table to define future structure of the table



### Admin Move Table - to avoid reorg pending (2|3)

```
[[db2inst1@db2_test ~]$ db2 "create table testschema.testtable (col1 int not null primary key, col2 int,col3 int,col4 char(50)) "
DB20000I The SQL command completed successfully.
[[db2inst1@db2_test ~]$ db2 "create index testschema.clus_idx on testschema.testtable (col2) cluster"
DB20000I The SQL command completed successfully.
[[db2inst1@db2_test ~]$ db2 "insert into testschema.testtable values (1,2,3,'FOUR')"
DB20000I The SQL command completed successfully.
[(reverse-i-search)`create': db2 "^Ceate index testschema.clus_idx on testschema.testtable (col2) cluster"
[[db2inst1@db2_test ~]$ db2 "create table testschema.testtable_new (col1 int not null primary key, col2 int,col4 char(50)) "
DB20000I The SQL command completed successfully.
```

[[db2inst1@db2\_test ~]\$ db2 "call SYSPROC.ADMIN\_MOVE\_TABLE('TESTSCHEMA','TESTTABLE','TESTTABLE\_NEW','','MOVE')";

KEY	VALUE
AUTHID CLEANUP_END CLEANUP_START COPY_END COPY_INDEXNAME COPY_INDEXSCHEMA COPY_OPTS COPY_START COPY_TOTAL_ROWS INDEXNAME INDEXSCHEMA INDEX_CREATION_TOTAL_TIME INIT_END INIT_START ORIGINAL_TBLSIZE PAR_COLDEF REPLAY_END REPLAY_END REPLAY_TOTAL_ROWS REPLAY_TOTAL_TIME STATUS SWAP_END SWAP_END SWAP_END SWAP_END	DB2INST1 2020-07-31-12.35.58.608997 2020-07-31-12.35.58.512877 2020-07-31-12.35.58.142783 CLUS_IDX TESTSCHEMA OVER_INDEX, ARRAY_INSERT, CLUSTER_OVER_INDEX 2020-07-31-12.35.58.114566 1 SQL200731123507230 SYSIBM 0 2020-07-31-12.35.58.037016 2020-07-31-12.35.57.702908 512 using a supplied target table so COLDEF could be different 2020-07-31-12.35.58.437970 2020-07-31-12.35.58.143856 0 0 COMPLETE 2020-07-31-12.35.58.468269 0
SWAP_START UTILITY_INVOCATION_ID VERSION	2020-07-31-12.35.58.460685 01000000B800000080000000000000000000000



#### Admin Move Table - to avoid reorg pending (3|3)

```
('TESTSCHEMA', 'TESTTABLE', ",",",",",",",", 'col1 int not null primary key, col2 int not null, col4 varchar(255)', ",'MOVE')";
```

```
[db2inst1@db2_test ~]$ db2 "call SYSPROC.ADMIN_MOVE_TABLE('TESTSCHEMA','TESTTABLE','','','','','','',''col1 int not null primary key,col2 int not null, col4 varchar(255)','','MOVE')";
 Result set 1
                                   VALUE
 AUTHID
                                   DB2INST1
 CLEANUP END
                                   2020-07-31-12.26.49.964842
 CLEANUP_START
                                   2020-07-31-12.26.49.854156
                                   2020-07-31-12.26.49.255086
 COPY END
 COPY_INDEXNAME
                                   CLUS IDX
 COPY_INDEXSCHEMA
                                   TESTSCHEMA
 COPY_OPTS
                                   OVER_INDEX, ARRAY_INSERT, CLUSTER_OVER_INDEX
 COPY_START
                                   2020-07-31-12.26.49.225587
 COPY_TOTAL_ROWS
 INDEXNAME
                                   SQL200731122629560
 INDEXSCHEMA
                                   SYSIBM
 INDEX_CREATION_TOTAL_TIME
 INIT_END
                                   2020-07-31-12.26.49.132874
 INIT_START
                                   2020-07-31-12.26.48.647045
 ORIGINAL TBLSIZE
 PAR_COLDEF
                                   col1 int not null primary key,col2 int not null, col4 varchar(255)
 REPLAY END
                                   2020-07-31-12.26.49.736263
                                   2020-07-31-12.26.49.255982
 REPLAY_START
 REPLAY_TOTAL_ROWS
 REPLAY TOTAL TIME
 STATUS
                                   COMPLETE
                                   2020-07-31-12.26.49.795169
 SWAP_END
 SWAP_RETRIES
 SWAP_START
                                   2020-07-31-12.26.49.783091
 UTILITY_INVOCATION_ID
                                   01000000AE0000000800000000000000000202007311226491340490000000
 VERSION
                                   11.05.0400
 26 record(s) selected.
 Return Status = 0
```



#### Performance Comparison - Backround (1 | 2)

- Project goal: verify reorg performance for the Db2 LUW database
- Test-System
  - AWS EC2 instance of the r5.12xlarge type
  - 48 vCPU
  - 384 GiB Memory
  - 8 EBS Volumes with 1TB and limited to 10.000 IOPS for db2 data
- Test-Tables (partitioned)
  - Tab1 1.201.894 pages (32k) → ~36GB | 142.572.100 Rows
  - Tab2 698.711 pages (32k) → ~21GB | 98.459.199 Rows
- Roughly gone through 150 test steps in total

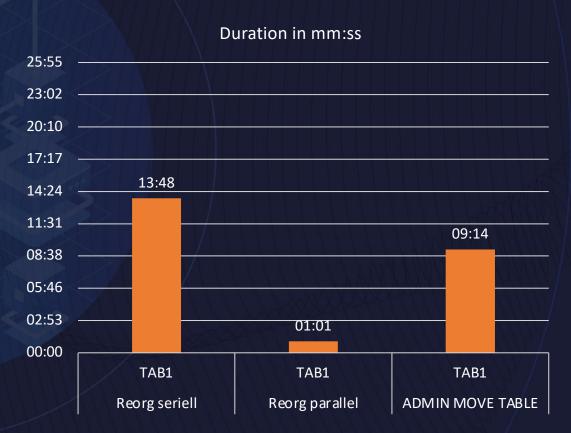


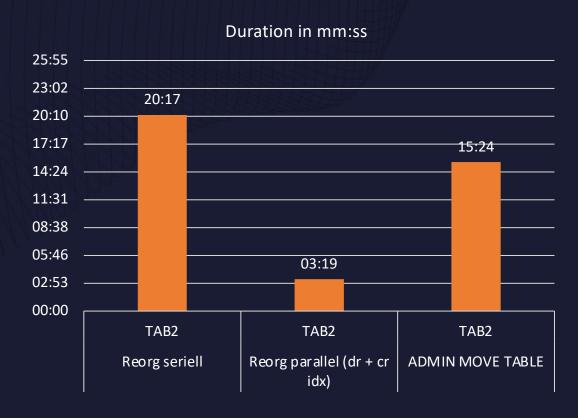
#### Performance Comparison - Backround (2 | 2)

- Recommended final solution has to be simple and out-of-the-box
- Data had the size of a test-system first, increased over time to see production like runtimes
- Both tested tables are the biggest in production, are having a clustered index and are partitioned
- Tab1 has no nonpartitioned index
- Tab2 has some nonpartitioned index ( > drop/create approach)
- Goal: Achive maximum ressource consumption with high parallelism
- Parallel: all 16 Partitions at once in parallel
- > Test: normal reorg vs. parallel (per partition) vs. admin move table



# Performance Comparison - default vs. parallel reorg vs. admin\_move\_table



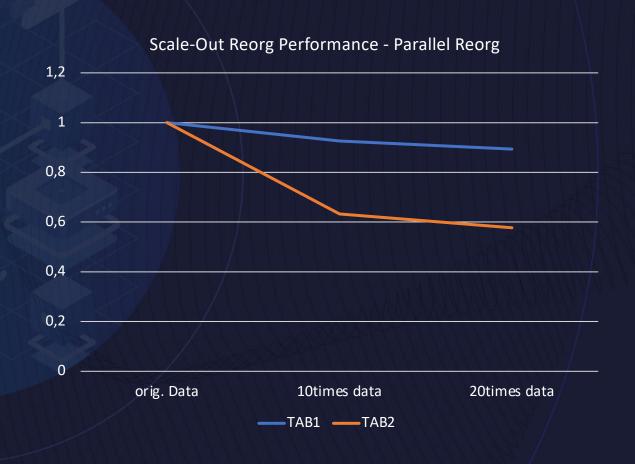


142Mio Rows, 36GB, all partitioned

98Mio Rows, 21GB, not all partitioned



#### Performance Comparison - scale-out parallel reorg



#### **Duration parallel reorg mm:ss**

	orig. Data	10times data	20times data
TAB1	01:01	09:25	18:10
TAB2	03:19	20:58	38:17

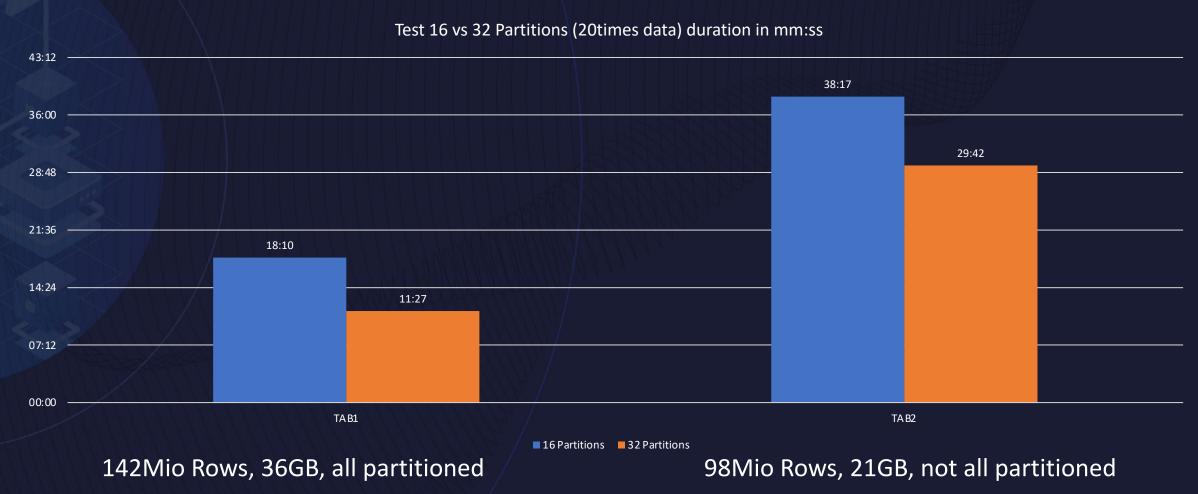
#### It gets faster when data increases!

TAB1: 142Mio Rows, 36GB, all partitioned

TAB2: 98Mio Rows, 21GB, not all partitioned



#### Performance Comparison - increase parallelism





#### Summary

- If you need fast reorg performance, use range partitioned tables (with a well balanced distribution key) and the reorg utility in parallel
- Reorg as many partitions in parallel as possible (rule of thumb when activity is low: Number of Cores -1, else: number of "free" cores)
- Avoid nonpartitioned indexes (or drop/recreate those)
- Use Admin Move Table for most of your every day reorgs (reorg recommended / improve performance / remove gaps)
- Avoid Reorg Pending, use admin move table instead!



#### Sources

https://www.sap.com/documents/2015/07/3080d083-5b7c-0010-82c7-eda71af511fa.html



**IDUG DB2 Tech Conference** 

Prague, Czech Republic | November 2011

# ADMIN\_MOVE\_TABLE() The New DBA Swiss Army Knife

Philip Nelson
Lloyds Banking Group / ScotDB Limited

Session Code: D1902



**IDUG DB2 Tech Conference** 

Prague, Czech Republic | November 2011

## Deep Dive in DB2 LUW offline table reorg and index reorg

Saeid Mohseni

Folksam Insurance Co.

Session Code: D11



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Don't forget to fill out a session evaluation!